A Distributed Triangulation Algorithm for Wireless Sensor Networks on 2D and 3D Surface

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OUTLINE

- Introduction
- Triangulation Algorithm
- Extension to 3D Surface
- Application and Simulation Results
- Conclusion



INTRODUCTION

- Wireless sensor network exhibits randomness
- Wireless sensor network topology is a graph
- Triangulation is a subgraph of the topology
- Triangulation mesh is very important for many applications of sensor networks: geometry-based routing, localization, coverage, segmentation, data segmentation



INTRODUCTION



(f) CDG (k=2): triangulated but coarse.

[2] R. Sarkar, X. Yin, J. Gao, F. Luo, and X. D. Gu, "Greedy Routing with Guaranteed Delivery Using Ricci Flows," in *Proc. of IPSN*, 2009.
[3] H. Zhou, S. Xia, M. Jin, and H. Wu, "Localized Algorithm for Precise Boundary Detection in 3D Wireless Networks," in Proc. of ICDCS, 2010.

INTRODUCTION

- We proposed a distributed triangulation algorithm
 - works for any arbitrary 2D sensor networks without communication model constraint, and we prove the correctness of the algorithm in 2D sensor networks
 - tolerates some measurement errors
 - also works for 3D open or closed surfaces



- Basic idea
 - Each triangulation mesh edge will associate with **two** triangles, except boundary edge
 - If we add extra edges into triangulation mesh, they will change this association
 - The association can be used to identify these extra edges.

- Definitions
 - node neighbor set (NNS), N_v(i)={h,k,l,m}
 - *edge neighbor set* (ENS), $N_e(e_{ij})=N_v(i)\cap N_v(j)=\{k,l,m\}$
 - refined edge neighbor set (RNS), $R_e(e_{ij}) = \{l,m\}$
 - $edge weight, W(e_{ij})=2, W(e_{ij})=2$
 - associated edge neight $(AENS), A(e_{ij}) = \{e_{il}, e_{im}, e_{jl}, e_{jl}\}$
 - equivalent edges, e_{ij}, e_{lm}
 - critical edge, ekj

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- Lemma 1: A subgraph of *G* is triangulated *if and only if* every edge of the subgraph has a weight of **two**.
- Given a graph G and a *triangulation subgraph T, extra edges are G-T*
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- Lemma 2: An *e*₀ *extra edge* can exist in *G*", only if it depends on at least two other *extra edges*.
- Lemma 3: An *e*₁ extra edge can exist in *G*", only if it at least depends on another extra edge.

(a) Loop chain.

(b) Non-loop chain. (c) Same head and tail.

- Step1: Initialization. each edge calculate *NNS*, *ENS*, *RENS*, *AENS*, and *edge weight*;
- Step2: Iterative edge removal according to **Theorems** 1 and 2;
- Step3: Removal of e_0 and e_1 chains.

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Application and Simulation Results

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(a) Triangulation under 10% distance errors.

(c) Triangulation under30% distance errors.

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Application and Simulation Results

Greedy routing based on Ricci flow in 2D networks

Application and Simulation Results

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COMPARISON OF STRETCH FACTOR IN GREEDY ROUTING.

	2D networks	3D networks
Triangulation by [2], [3]	2.065	5.417
Proposed Triangulation	1.214	1.091

Distribution of stretch factor in greedy routing

(b) 3D networks

CONCLUSION

- In summary, we have proposed a distributed algorithm that can triangulate an arbitrary sensor network without position information.
- The algorithm can achieve the finest triangulation and tolerate distance measurement errors.
- We have proven its correctness in 2D and extended it to 3D surface.
- A range of geometry-based network algorithms can benefit from the proposed triangulation.

